

Complete Addition Formulæ for Elliptic Curves

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Consider the elliptic curve E over a field \mathbb{K} (with $\text{Char } \mathbb{K} \neq 2, 3$) given by a Weierstraß equation

$$Y^2Z = X^3 + aXZ^2 + bZ^3$$

where $a, b \in \mathbb{K}$ are constants, $4a^3 + 27b^2 \neq 0$. The set $E(\mathbb{K})$ of points $(X : Y : Z) \in \mathbb{P}^2(\mathbb{K})$ forms an abelian group under the chord-and-tangent rule, with neutral element $\mathbf{O} = (0 : 1 : 0)$. The addition law is written additively. The negative of a point $\mathbf{P}_1 = (X_1 : Y_1 : Z_1)$ is $(X_1 : -Y_1 : Z_1)$. Given two points $\mathbf{P}_1 = (X_1 : Y_1 : Z_1)$ and $\mathbf{P}_2 = (X_2 : Y_2 : Z_2)$, we let $\mathbf{P}_3 := \mathbf{P}_1 + \mathbf{P}_2 = (X_3 : Y_3 : Z_3)$.

ADDITION ALGORITHM. The usual algorithm (*e.g.*, see [6, Chapter III, §2]) for adding two points on $E(\mathbb{K})$ distinguishes several cases:

1. If $\mathbf{P}_1 = \mathbf{O}$ then $\mathbf{P}_3 = (X_2 : Y_2 : Z_2)$;
2. If $\mathbf{P}_2 = \mathbf{O}$ then $\mathbf{P}_3 = (X_1 : Y_1 : Z_1)$;
3. If $\mathbf{P}_1 = -\mathbf{P}_2$ then $\mathbf{P}_3 = (0 : 1 : 0)$;
4. If $\mathbf{P}_1 \neq \pm\mathbf{P}_2$ and $\mathbf{P}_1, \mathbf{P}_2 \neq \mathbf{O}$ then $\mathbf{P}_3 = (X_3 : Y_3 : Z_3)$ where

$$\begin{cases} X_3 = (X_1Z_2 - X_2Z_1)X'_3 \\ Y_3 = (Y_1Z_2 - Y_2Z_1)[(X_1Z_2 - X_2Z_1)^2 X_1Z_2 - X'_3] - (X_1Z_2 - X_2Z_1)^3 Y_1Z_2 \\ Z_3 = (X_1Z_2 - X_2Z_1)^3 Z_1Z_2 \end{cases} \quad (1)$$

with $X'_3 = (Y_1Z_2 - Y_2Z_1)^2 Z_1Z_2 + (X_1Z_2 - X_2Z_1)^3 - 2(X_1Z_2 - X_2Z_1)^2 X_1Z_2$.

5. If $\mathbf{P}_1 = \mathbf{P}_2 \neq \mathbf{O}$ then $\mathbf{P}_3 = (X_3 : Y_3 : Z_3)$ where

$$\begin{cases} X_3 = 2Y_1Z_1[(3X_1^2 + aZ_1^2)^2 - 8X_1Y_1^2Z_1] \\ Y_3 = (3X_1^2 + aZ_1^2)[12X_1Y_1^2Z_1 - (3X_1^2 + aZ_1^2)^2] - 2(2Y_1^2Z_1)^2 \\ Z_3 = (2Y_1Z_1)^3 \end{cases} \quad (2)$$

Borrowing the notation of [1], \mathbf{M} , \mathbf{S} , and \mathbf{add} will respectively stand for the cost of a field multiplication, a field squaring, and a field addition (in \mathbb{K}), and $\mathbf{*c}$ will stand for the cost of the multiplication by some given constant $c \in \mathbb{K}$.

A careful operation count shows that the addition operation (1) costs $\underline{12\mathbf{M} + 2\mathbf{S} + 6\mathbf{add} + 1*2}$ and that the doubling operation (2) costs $\underline{5\mathbf{M} + 6\mathbf{S} + 1*a + 7\mathbf{add} + 3*2 + 1*3}$ [1,3].

A COMPLETE ADDITION LAW. It is worth noting that formulæ (1) and (2) are not valid for the point at infinity \mathbf{O} . We present below a formula that is valid for \mathbf{O} , as well as for the case $\mathbf{P}_1 = \mathbf{P}_2$. It is adapted from Formula III in [4, Section 3] and optimized.

Let $\mathbf{P}_1 = (X_1 : Y_1 : Z_1)$ and $\mathbf{P}_2 = (X_2 : Y_2 : Z_2)$. We assume that $\mathbf{P}_1 - \mathbf{P}_2$ is not a finite¹ point of order 2. Then $\mathbf{P}_3 := \mathbf{P}_1 + \mathbf{P}_2 = (X_3 : Y_3 : Z_3)$ is given by

$$\begin{cases} X_3 = (X_1Y_2 + X_2Y_1)[Y_1Y_2 - 3bZ_1Z_2 - a(X_1Z_2 + X_2Z_1)] - \\ \quad (Y_1Z_2 + Y_2Z_1)[a(X_1X_2 - aZ_1Z_2) + 3b(X_1Z_2 + X_2Z_1)] \\ Y_3 = (Y_1Y_2 + 3bZ_1Z_2)(Y_1Y_2 - 3bZ_1Z_2) + a(X_1X_2 - aZ_1Z_2)(3X_1X_2 + aZ_1Z_2) + \\ \quad (X_1Z_2 + X_2Z_1)[3b(X_1X_2 - aZ_1Z_2) - a^2(X_1Z_2 + X_2Z_1)] \\ Z_3 = (Y_1Z_2 + Y_2Z_1)[Y_1Y_2 + 3bZ_1Z_2 + a(X_1Z_2 + X_2Z_1)] + (X_1Y_2 + X_2Y_1)(3X_1X_2 + aZ_1Z_2) \end{cases} \quad (3)$$

Remark 1. If $\mathbf{P}_1 - \mathbf{P}_2 = (\xi : 0 : 1)$ for some $\xi \in \mathbb{K}$ (i.e., it is a finite point of order 2) then $\mathbf{P}_3 = \mathbf{P}_1 + \mathbf{P}_2 = (X_3 : Y_3 : Z_3)$ is given by the usual addition algorithm; namely, \mathbf{P}_3 is given by Eq. (1) if $\mathbf{P}_1, \mathbf{P}_2 \neq \mathbf{O}$, and by $\mathbf{P}_3 = (\xi : 0 : 1)$ if \mathbf{P}_1 or $\mathbf{P}_2 = \mathbf{O}$. As demonstrated in [2], the case $\mathbf{P}_1 - \mathbf{P}_2 = (\xi : 0 : 1)$ for some $\xi \in \mathbb{K}$ can also be handled by a single addition formula.

Remark 2. Reference [4] produces two other addition formulas that can handle \mathbf{O} . However, they do not fit our needs. The same formulas, extended to the long Weierstraß equations, can be found in [5,2].

DETAILED ALGORITHM AND COMPLEXITY.

$$\begin{aligned} X1X2 &= X1*X2 \\ Y1Y2 &= Y1*Y2 \\ Z1Z2 &= Z1*Z2 \\ u &= (X1+Y1)*(X2+Y2) - X1X2 - Y1Y2 \\ v &= (X1+Z1)*(X2+Z2) - X1X2 - Z1Z2 \\ w &= (Y1+Z1)*(Y2+Z2) - Y1Y2 - Z1Z2 \\ Za &= a*Z1Z2 \\ Zb &= 3b*Z1Z2 \\ va &= a*v \\ P &= X1X2 + Za \\ Q &= X1X2 - Za \\ Qa &= a*Q \\ M &= 2*X1X2 + P \\ R &= Y1Y2 + Zb \\ S &= Y1Y2 - Zb \\ X3 &= u*(S - va) - w*(Qa + 3b*v) \\ Y3 &= R*S + Qa*M + v*(3b*Q - a*va) \\ Z3 &= w*(R + va) + u*M \end{aligned}$$

Cost: $\underline{13M + 4*a + 3*3b + 25add + 1*2}$.

References

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¹ By finite point, we mean a point with a Z -coordinate different from 0.

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